

Compressed sensing with convex optimization

Note Title

2015-10-31

Recall

Thm Master theorem
 Let $A \in \mathbb{R}^{m \times n}$ have independent entries satisfying:
 $\mathbb{E} A_{ij} = 0$, $\mathbb{E} A_{ij}^2 = 1$, $\|A_{ij}\|_{\ell_2} \leq 10$.
 Let $0 < \epsilon < \mathbb{R}^n$. Then

$$\mathbb{E} \sup_{x \in K} \|Ax\|_2 - \sqrt{m} \|x\|_2 \leq C w(\epsilon)$$

$\Rightarrow P(\sup_{x \in K} \|Ax\|_2 - \sqrt{m} \|x\|_2 \leq C w(\epsilon)) \geq 99\%$

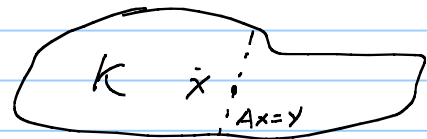
Markov
 \neq

Can be any const.

Application: signal recovery in structured linear model:

$$y = Ax, \quad x \in K \subseteq \mathbb{R}^n$$

sub-Gaussian as in "master thm"



Recover, or estimate x , by solving program:

$$\text{Find } x' \in K \text{ s.t. } y = Ax' \quad (1)$$

Let \hat{x} be the solution.

Accuracy of \hat{x} ?

Note: $\hat{x}, x \in K \quad A\hat{x} = Ax = y$.

Let $h = \hat{x} - x$. Then $h \in K - K$, $h \in \text{null}(A)$

$\|h\|_2$ is controlled via the low m^* estimate, which follows from the above thm.

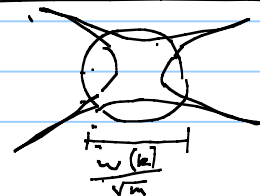
Thm (Low m^* estimate)

Under the conditions of theorem above

$$\text{diam}(\text{null}(A) \cap K) \leq \frac{C w(K)}{\sqrt{m}}, \quad \text{w.p. } > 99\%$$

Euclidean diameter

random subspace w/ codimension m .



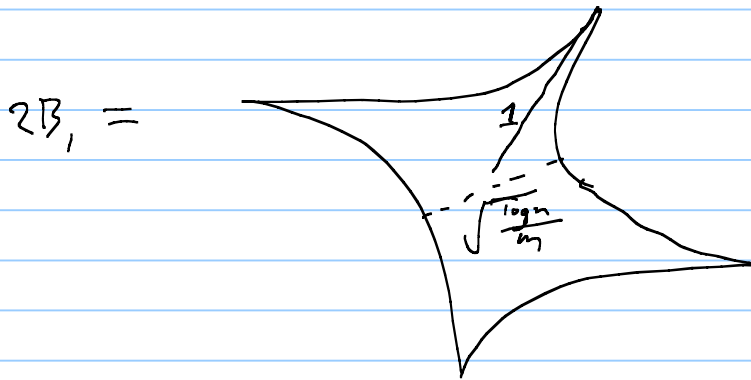
Thus, w.p. $> 99\%$,

$$\|h\|_2 \leq C \frac{w(K-K)}{\sqrt{m}}$$

Ex) $K = B_1$,

$$w(K-K) = w(2B_1) = \mathbb{E} \sup_{x \in 2B_1} \langle x, g \rangle = \mathbb{E} \|21g\|_2 \leq C \sqrt{\log n}$$

$$\Rightarrow \|h\|_2 \leq C \sqrt{\frac{\log n}{m}}$$



Signal processing interpretation:

As soon as $m \gg \sqrt{\log n}$, $sola$ becomes very accurate.

Geometric interpretation: Intersection of B_1 w/ a random hyperplane of codimension $m \gg \sqrt{\log n}$ has much smaller Euclidean diameter than that of B_1 .

However, the approach above can be considerably improved.

Q: What if K is unbounded $\Rightarrow w(K) = \infty$?

Q: What if the program (1) is computationally intractible?

Both problems occur when $K_S = \{x \in \mathbb{R}^n : \|x\|_0 \leq S\}$.

Solution: convexify and examine local properties of feasible set.

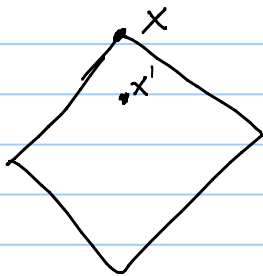
Motivating example:

$$y = Ax, \quad x \in K_S$$

Not convex. Replace w/ B_1^n .

Let \hat{x} be a soln to the convex program:

$$\text{Find } x' \text{ s.t. } Ax = Ax', \quad \|x'\|_1 \leq \|x\|_1,$$

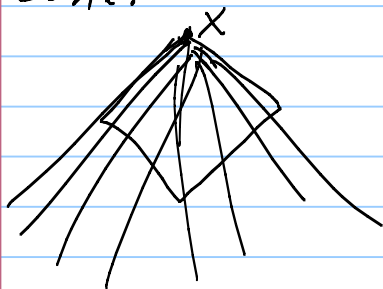


Set $h := \hat{x} - x$. Note:

(1) $h \in N(A)$

(2) $\|x+h\|_1 \leq \|x\|_1$

Locally, the second constraint looks like a cone.



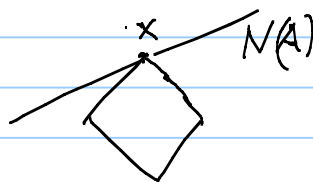
Def (Tangent cone):

$$D(K, x) := \{ \tau(v-x) : v \in K, \tau \geq 0 \}$$

Observe:

a) $h \in D(K, x)$

b) If $D(K, x) \cap N(A) = \{0\}$ then $h=0$.



i.e. $N(A)$ "escapes" the tangent cone.

The prob. of this good event is captured by the following theorem:

Thm (Gordon's "escape through the mesh" thm) —

Let A be sub-Gaussian as in the "master theorem". Let D be a cone.

Let

$$\epsilon = C \frac{w(D \cap B_2^n)}{\sqrt{m}}.$$

Then w.p. $\geq 99\%$, $\forall x \in D$,

$$A) (1-\epsilon) \|x\|_2 \leq \frac{\|Ax\|_2}{\sqrt{m}} \leq (1+\epsilon) \|x\|_2.$$

In particular, if $m \geq C w(D \cap B_2^n)$, then

$$B) n(A) \cap D = \{0\} \quad \text{w.p. } \geq 99\%.$$

A) "A is well conditioned when restricted to D if $m \geq C w(D \cap B_2^n)$."

Proof B) follows from the lower bound of A).

Exercise: Prove A) based on "master thm".

Return to example w/ $x \in K_S$, $h \in D(\|x\|; B_1^n, x)$.

One can show that

$$w(D(\|x\|; B_1^n, x) \cap B_2^n) \leq C s \log\left(\frac{cn}{s}\right).$$

Thus, "escape through mesh" thm implies that $\hat{x} = x$ w.p. $\geq 99\%$ if $m \geq C s \log\left(\frac{cn}{s}\right)$.

Astonished researcher (2004): "x is recovered exactly! By a convex program! w/ hardly more measurements than would be needed for ℓ_0 minimization!"

General analysis using tangent cone:

Let \hat{x} be a soln to $\begin{matrix} \text{convex set} \\ \text{containing } x. \end{matrix}$

Find x' s.t. $Ax' = y, x' \in K$ (2)

Corollary If $m \geq c w(D(K, x) \cap B_c^n)^2$, then
w.p. $> 99\%$, $\tilde{x} = x$.

Proof: Same steps as in example above.

Remarks

- (2) is often computationally tractable.
- K is not necessarily the signal structure, rather it is a convex surrogate. Given a certain signal structure, determining what K to use is a question of interest.
- In some cases $w(K, x)^2$ can be tiny. This is what allows $m \ll n$, i.e., "dimension reduction".

